

1 Definitions

Définition 1. natural transformation

Définition 2. Let $\mathcal{C}$ be a category.

A monad on the category $\mathcal{C}$ is (T, μ, η) where $T : \mathcal{C} \rightarrow \mathcal{C}$ is an endofunctor and where $\mu : TT \rightarrow T$ and $\eta : \text{id}_{\mathcal{C}} \rightarrow T$ are natural transformations such that the following diagrams commute :

$$\begin{array}{ccc} T \circ T \circ T & \xrightarrow{T \circ \mu} & T \circ T \\ \downarrow \mu \circ T & & \downarrow \mu \\ T \circ T & \xrightarrow{\mu} & T \end{array} \quad \begin{array}{ccc} T & \xrightarrow{\eta \circ T} & T \circ T \\ \downarrow T \circ \eta & \searrow id & \downarrow \mu \\ T \circ T & \xrightarrow{\mu} & T \end{array}$$

Définition 3. A monoidal category is $(\mathcal{C}, \otimes, I, \lambda, \rho, \alpha)$ where

- $\mathcal{C}$ is a category,
- $\otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ is a bifunctor of $\mathcal{C}$,
- I is an object of $\mathcal{C}$,
- λ and ρ are invertible natural transformations given for every objet A of $\mathcal{C}$ by

$$\lambda_A : I \otimes A \rightarrow A \quad \rho_A : A \otimes I \rightarrow A,$$

- α is an invertible natural transformation given for every objects A, B and C of $\mathcal{C}$ by

$$\alpha_{A,B,C} : (A \otimes B) \otimes C \rightarrow A \otimes (B \otimes C)$$

such that the following diagrams commute :

$$\begin{array}{ccc} ((A \otimes B) \otimes C) \otimes D & \xrightarrow{\alpha_{A,B,C \otimes D}} & (A \otimes (B \otimes C)) \otimes D \xrightarrow{\alpha_{A,B \otimes C,D}} A \otimes ((B \otimes C) \otimes D) \\ \downarrow \alpha_{A \otimes B,C,D} & & \downarrow A \otimes \alpha_{B,C,D} \\ (A \otimes B) \otimes (C \otimes D) & \xrightarrow{\alpha_{A,B,C \otimes D}} & A \otimes (B \otimes (C \otimes D)) \end{array}$$

$$\begin{array}{ccc} (A \otimes I) \otimes B & \xrightarrow{\alpha_{A,I,B}} & A \otimes (I \otimes B) \\ \searrow \rho_{A \otimes B} & & \swarrow A \otimes \lambda_B \\ & A \otimes B & \end{array}$$

Définition 4. A monoidal strict category $(\mathcal{C}, \otimes, I)$ is a monoidal category $(\mathcal{C}, \otimes, I, \text{id}_{\mathcal{C}}, \text{id}_{\mathcal{C}}, \text{id}_{\mathcal{C}})$.

monoid in $(Cat, \times, 1)$

Définition 5. A strong monad (T, η, μ, t) on a monoidal category $(\mathcal{C}, \otimes, I, \lambda, \rho, \alpha)$ is a monad (T, η, μ) on the category $\mathcal{C}$ equipped with a natural transformation t called right-force, defined by

$$t_{A,B} : A \otimes TB \rightarrow T(A \otimes B)$$

such that the following diagrams commute for any objects A, B and C of $\mathcal{C}$:

$$\begin{array}{ccc} I \otimes TA & \xrightarrow{t_{I,A}} & T(I \otimes A) \\ & \searrow \lambda_{TA} & \downarrow T\lambda_A \\ & & TA \end{array} \quad \begin{array}{ccc} A \otimes B & \xrightarrow{A \otimes \eta_B} & A \otimes TBt_{A,B} \\ & \searrow \eta_{A \otimes B} & \downarrow \\ & & T(A \otimes B) \end{array}$$

$$\begin{array}{ccc} (A \otimes B) \otimes TC & \xrightarrow{t_{A \otimes B, C}} & T((A \otimes B) \otimes C) \\ \alpha_{A,B,TC} \downarrow & & \downarrow T\alpha_{A,B,C} \\ A \otimes (B \otimes C) & \xrightarrow{A \otimes t_{B,C}} A \otimes T(B \otimes C) \xrightarrow{t_{A,B \otimes C}} & T(A \otimes (B \otimes C)) \end{array}$$

$$\begin{array}{ccccc} A \otimes TTB & \xrightarrow{t_{A,TB}} & T(A \otimes TB) & \xrightarrow{Tt_{A,B}} & TT(A \otimes B) \\ A \otimes \mu_B \downarrow & & & & \downarrow \mu_{A,B} \\ A \otimes TB & \xrightarrow{t_{A,B}} & T(A \otimes B) & & \end{array}$$

The notion of left-force on a monad is defined similarly.

Définition 6. A symmetric monoidal category $(\mathcal{C}, \otimes, I, \lambda, \rho, \alpha, \gamma)$ is a monoidal category $(\mathcal{C}, \otimes, I, \lambda, \rho, \alpha)$ equipped with a natural transformation γ defined for all objects A and B of $\mathcal{C}$ by

$$\gamma_{A,B} : A \otimes B \rightarrow B \otimes A$$

such that the following diagrams commute

$$\begin{array}{ccc} B \otimes A & \xrightarrow{\gamma_{B,A}} & A \otimes B \\ & \searrow id_{B \otimes A} & \downarrow \gamma_{A,B} \\ & & B \otimes A \end{array} \quad \begin{array}{ccccc} (A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\gamma_{A,B \otimes C}} & (B \otimes C) \otimes A \\ \gamma_{A,B \otimes C} \downarrow & & & & \downarrow \alpha_{B,C,A} \\ (B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{B \otimes \gamma_{A,C}} & B \otimes (C \otimes A) \end{array}$$

Lemme 7. Let $(\mathcal{C}, \otimes, I, \lambda, \rho, \alpha, \gamma)$ a symmetric monoidal category with a monad (T, η, μ) .

Every left-force on the monad induces a right-force on the monad. Conversely, every right-force on the monad induces a left-force on the monad.

Proof. Let t be a left-force. The natural transformation t induces a natural transformation s defined by :

$$s_{A,B} = T\gamma_{A,B} \circ t_{B,A} \circ \gamma_{TA,B} : TA \otimes B \rightarrow T(A \otimes B)$$

Let us now show that s is indeed a right-force.

$$\begin{array}{ccccc} TA \otimes I & \longrightarrow & I \otimes TA & & T(I \otimes A) & & T(A \otimes I) \\ & & \searrow (1) & \searrow & & & \\ & & & & TA & & \end{array}$$

□

2 Forests and monads

Définition 8. A forest is given by a set V and a mapping

$$p : V \rightarrow V \uplus \{\perp\}$$

such that